

# CS 465

# Computer Security

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MAC: Message Authentication Code

# What Assurances are Provided by Symmetric Encryption?

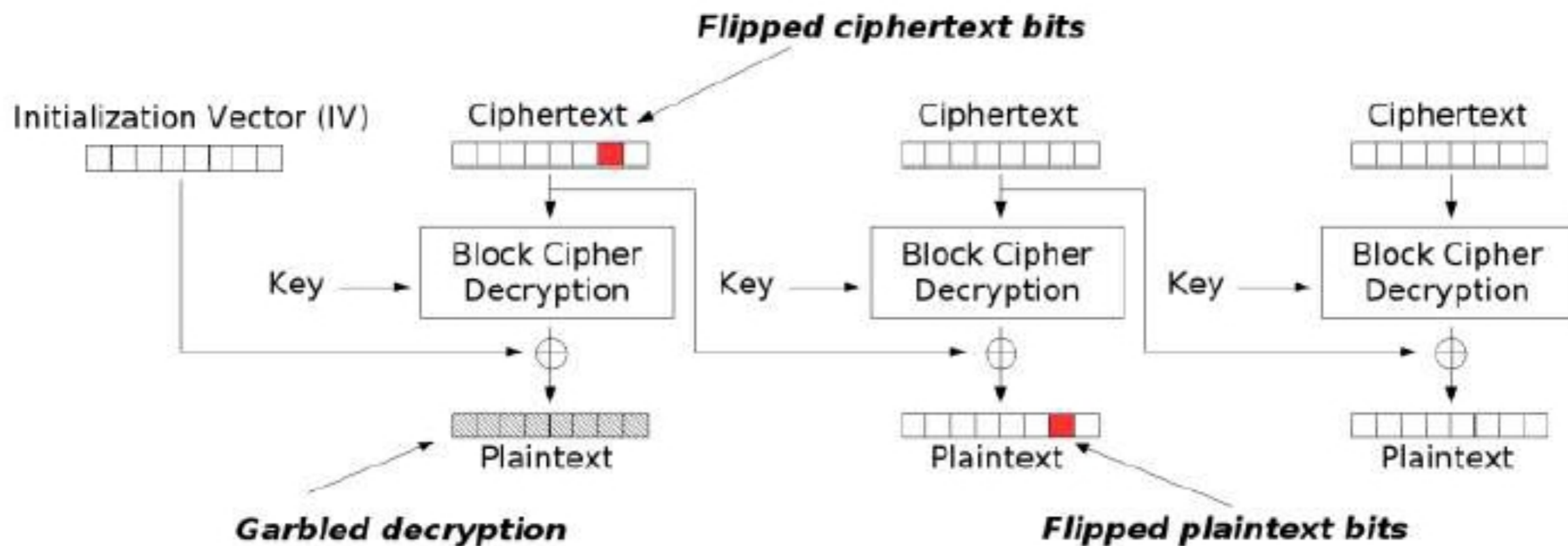
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Assume CTR or CBC mode

- Authentication?
- Confidentiality?
- Integrity?
- Non-repudiation?

# Bit Flipping Attacks (Block Cipher)

## Modification attacks on CBC



Modification attack on CBC

# Bit Flipping Attacks (Stream Cipher)

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- Plaintext:

- ACCT\_NO:123-45-6789 ADD:100

- Ciphertext:

- 15b1206b7efa68b9 89 c87357507e3a27a138ca dc b2a1bb f8 eebee5

# Goals of Message Authentication

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- Assure that the message has not been altered
- Assure the source of the message is authentic

# Message Authentication: Ciphertext vs. Plaintext

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- Authentication of encrypted messages
  - Include an error-detection code in plaintext message
  - Attach a key-based error-detection code to an encrypted message
  - Attach a TAG – remember the newer AEAD modes
- Authentication of plaintext messages
  - Authentication without confidentiality
  - Attach a key-based error-detection code to plaintext message

# Message Authentication Code (MAC)

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Dear BYU,  
  
Thank you so  
much for an  
awesome  
computer security  
course.

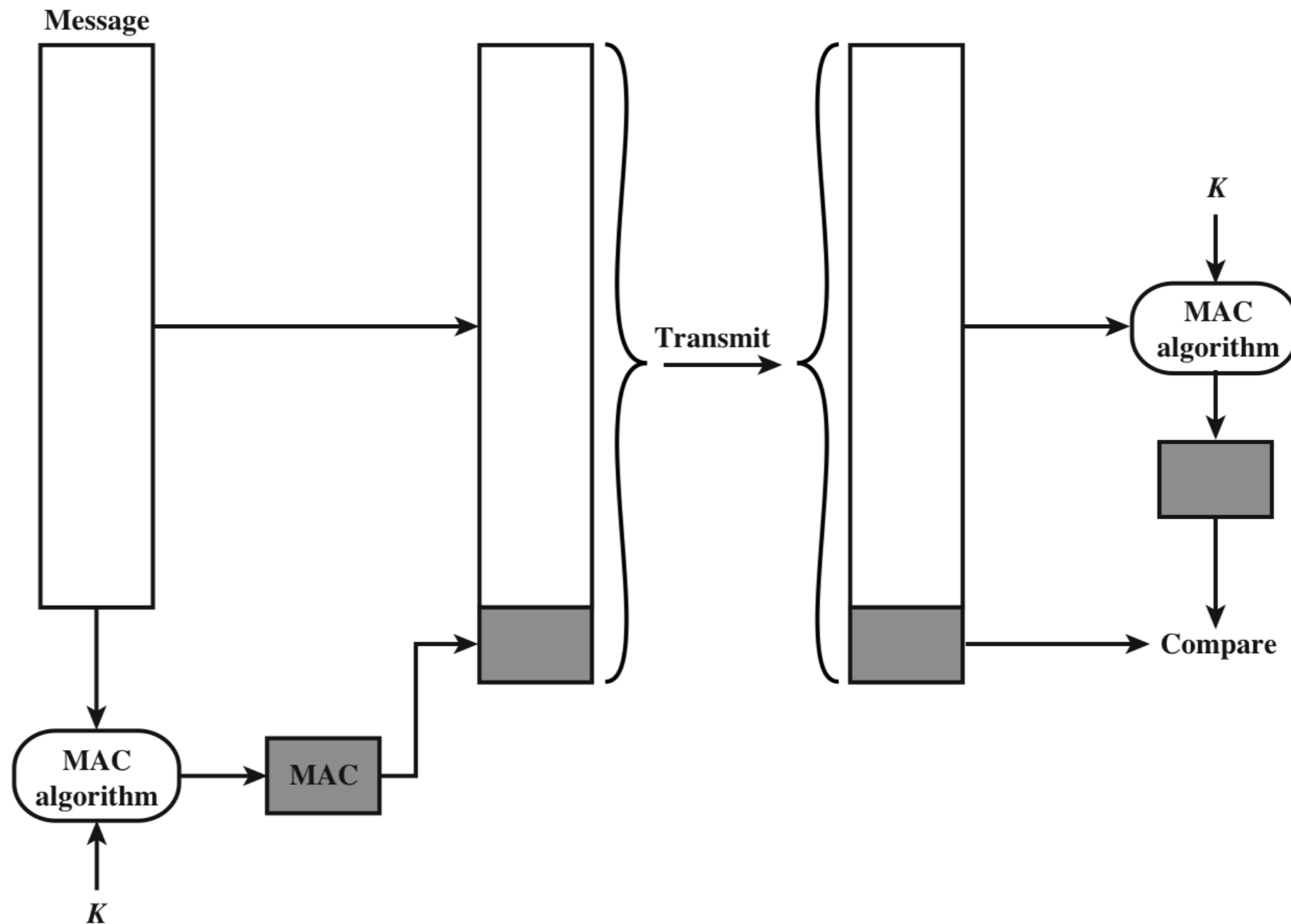
Sincerely,  
Emma



MAC Algorithm

This message  
really is from  
me and hasn't  
been modified

# Message Authentication Code (MAC)



**Figure 3.1** Message Authentication Using a Message Authentication Code (MAC)



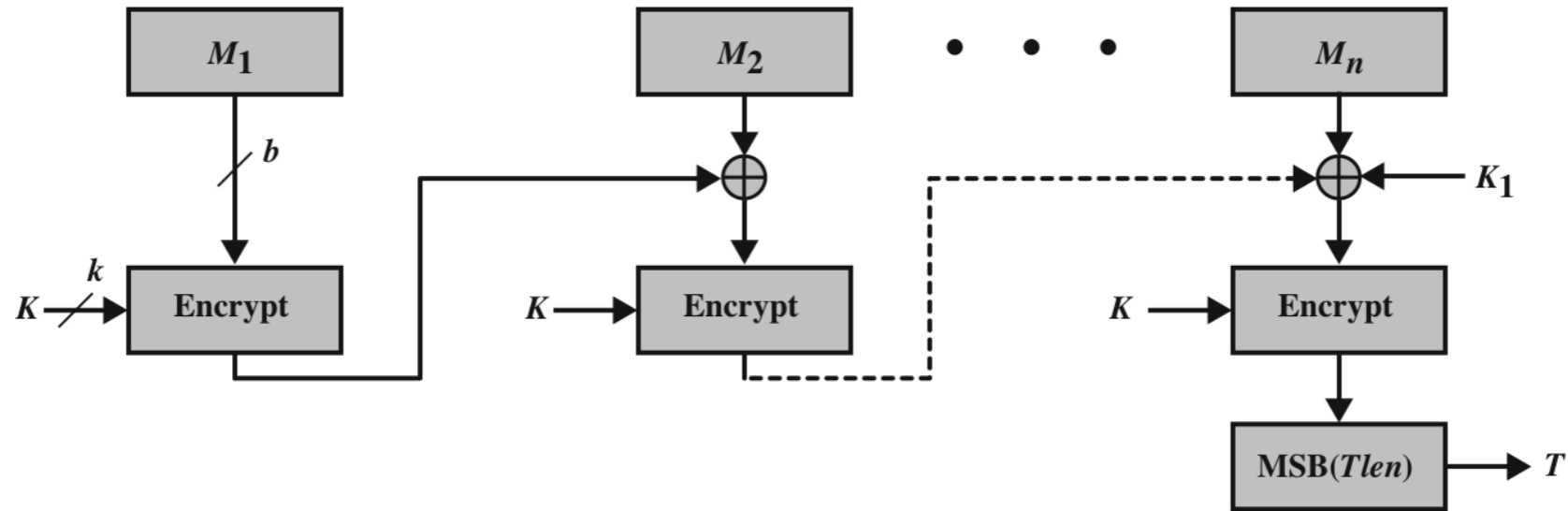
# Three Ways to Implement a MAC

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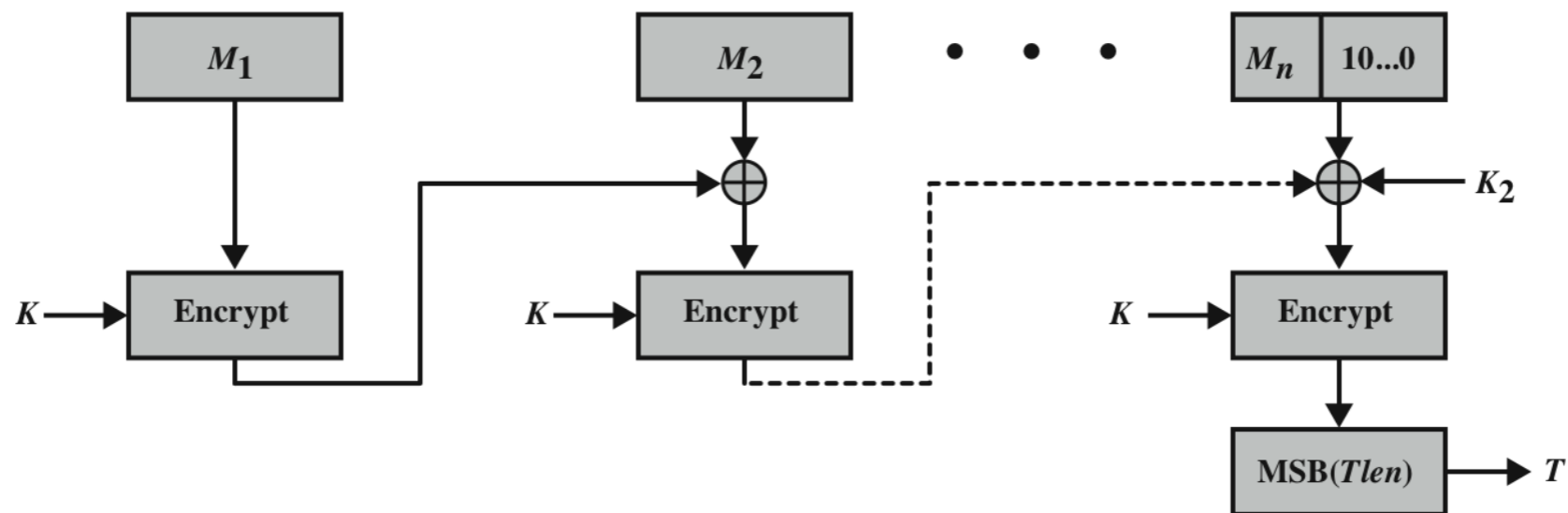
## 1. CBC-MAC

- Use CBC mode and a block cipher — fixed length messages only
- OMAC — for variable length messages

# OMAC1 (also called CMAC)



(a) Message length is integer multiple of block size

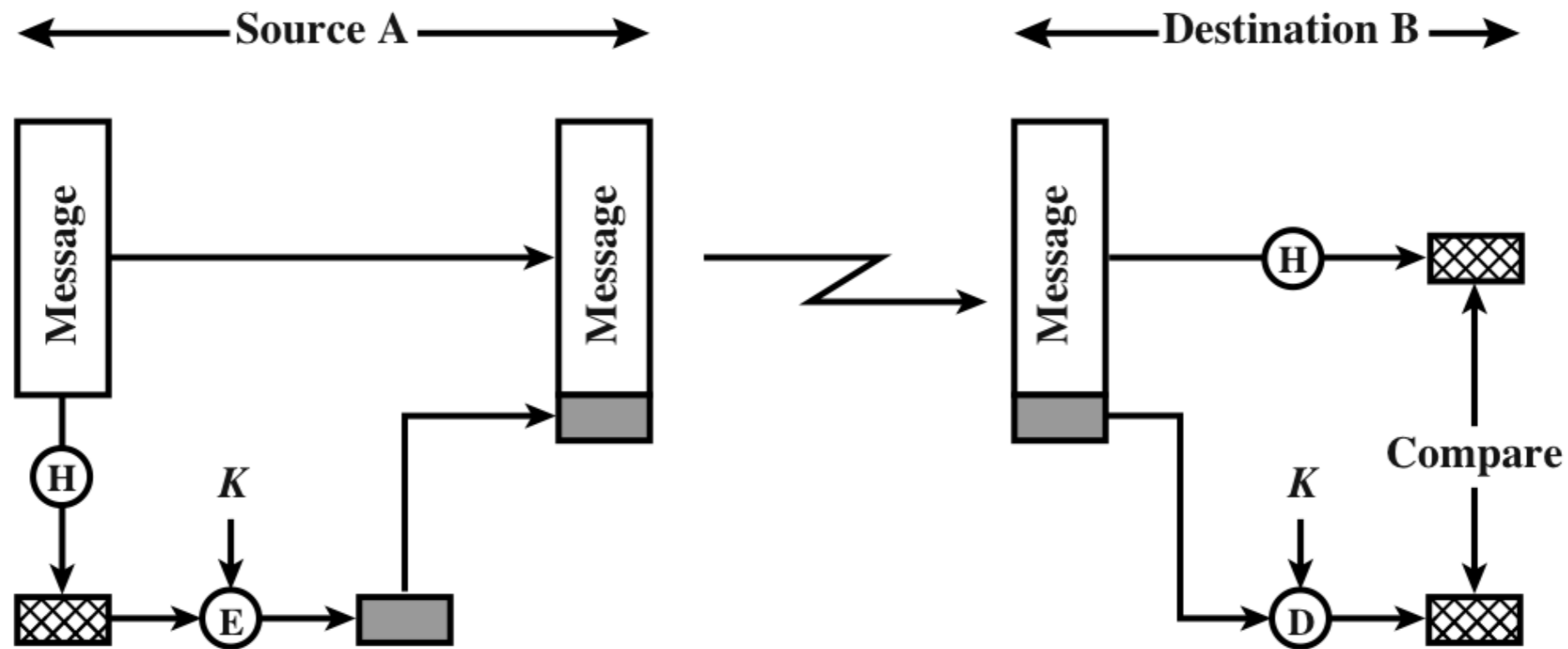


(b) Message length is not integer multiple of block size

Figure 12.12 Cipher-Based Message Authentication Code (CMAC)

# Three Ways to Implement a MAC

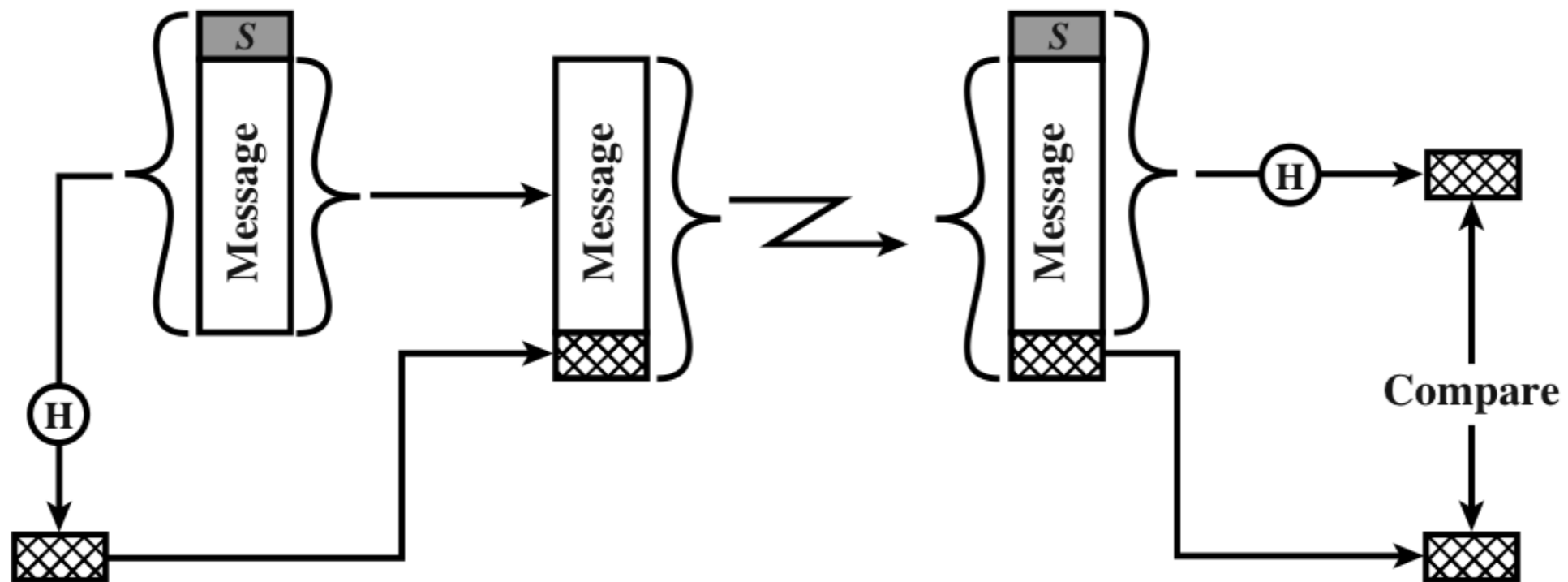
## 2. Hash the message and encrypt the digest



(a) Using conventional encryption

# Three Ways to Implement a MAC

## 3. Hash the message along with a shared key



(c) Using secret value

# Design Flaw!

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- Cryptographers recommend against this kind of MAC using modern hash functions
- Vulnerable to a message extension attack
- Vulnerability comes from an interactive implementation technique known as the Merkle-Damgård construction
- Hash functions that are vulnerable when used this way: MD5, SHA1, SHA2



# HMAC

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- Because of the message extension attack vulnerability, the government standard HMAC algorithm guards against this threat
  - FIPS 198
  - RFC 2104

$$\text{HMAC}(K, m) = H\left((K' \oplus \text{opad}) \parallel H((K' \oplus \text{ipad}) \parallel m)\right)$$

- $K' = H(K)$  if  $K$  is larger than the block size, otherwise  $K$
- $\text{opad} = 0x5c5c5c\dots5c5c$ , one-block-long constant
- $\text{ipad} = 0x363636\dots3636$ , one-block-long constant
- $IV$  is fixed, as with SHA-2 and other hash functions

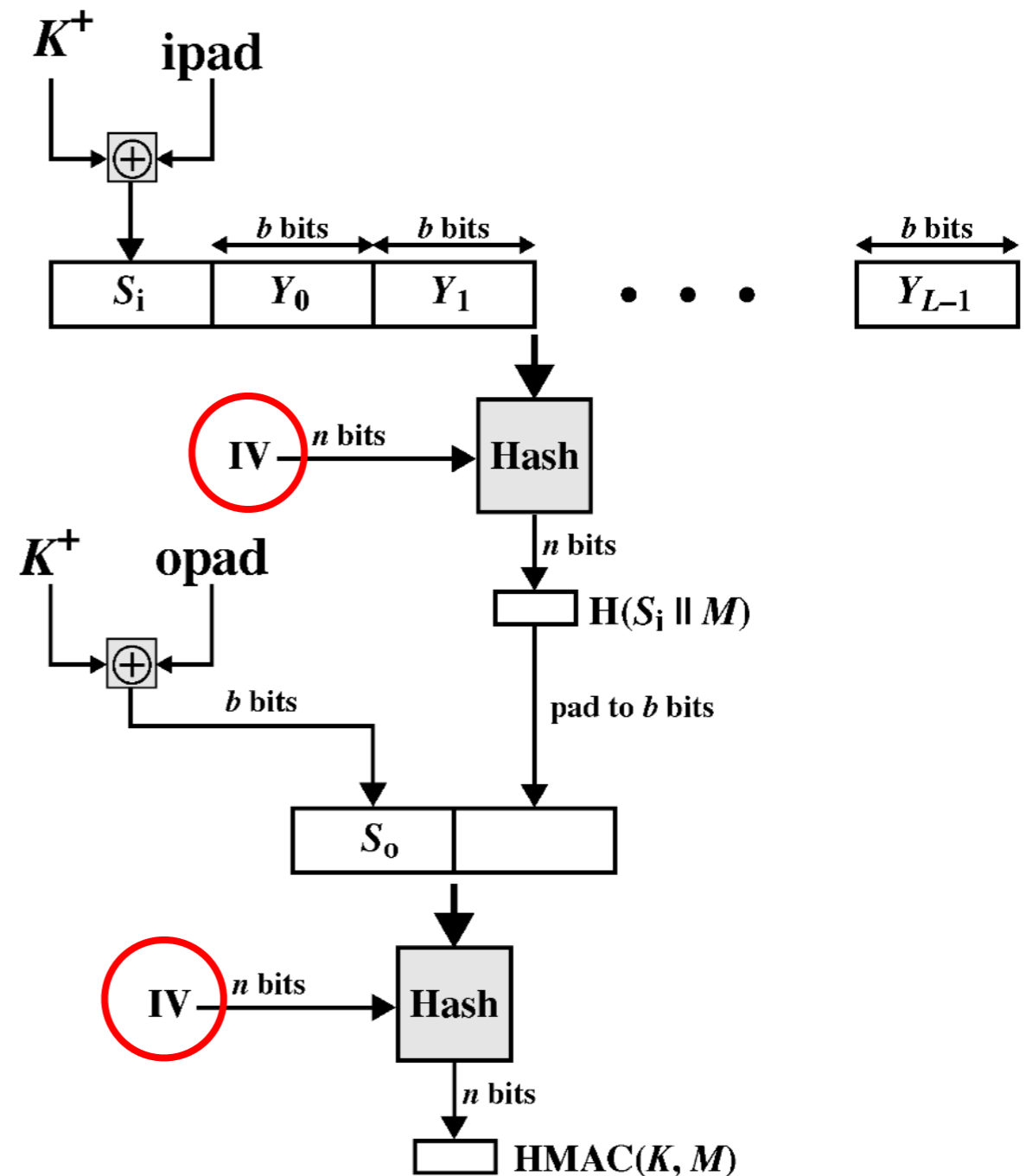


Figure 3.6 HMAC Structure

# HMAC = Hashed MAC

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- Wrong way:  $H(\text{secret} || m)$
- Right way :  $H((K' \oplus \text{opad}) || H((K' \oplus \text{ipad}) || m))$



# Recommendation

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- If you need just a MAC, use HMAC
- If you need encryption and a MAC, use AEAD
- See <https://blog.cryptographyengineering.com/2013/02/15/why-i-hate-cbc-mac/>

# MAC ATTACK

# Let's examine the message extension attack...

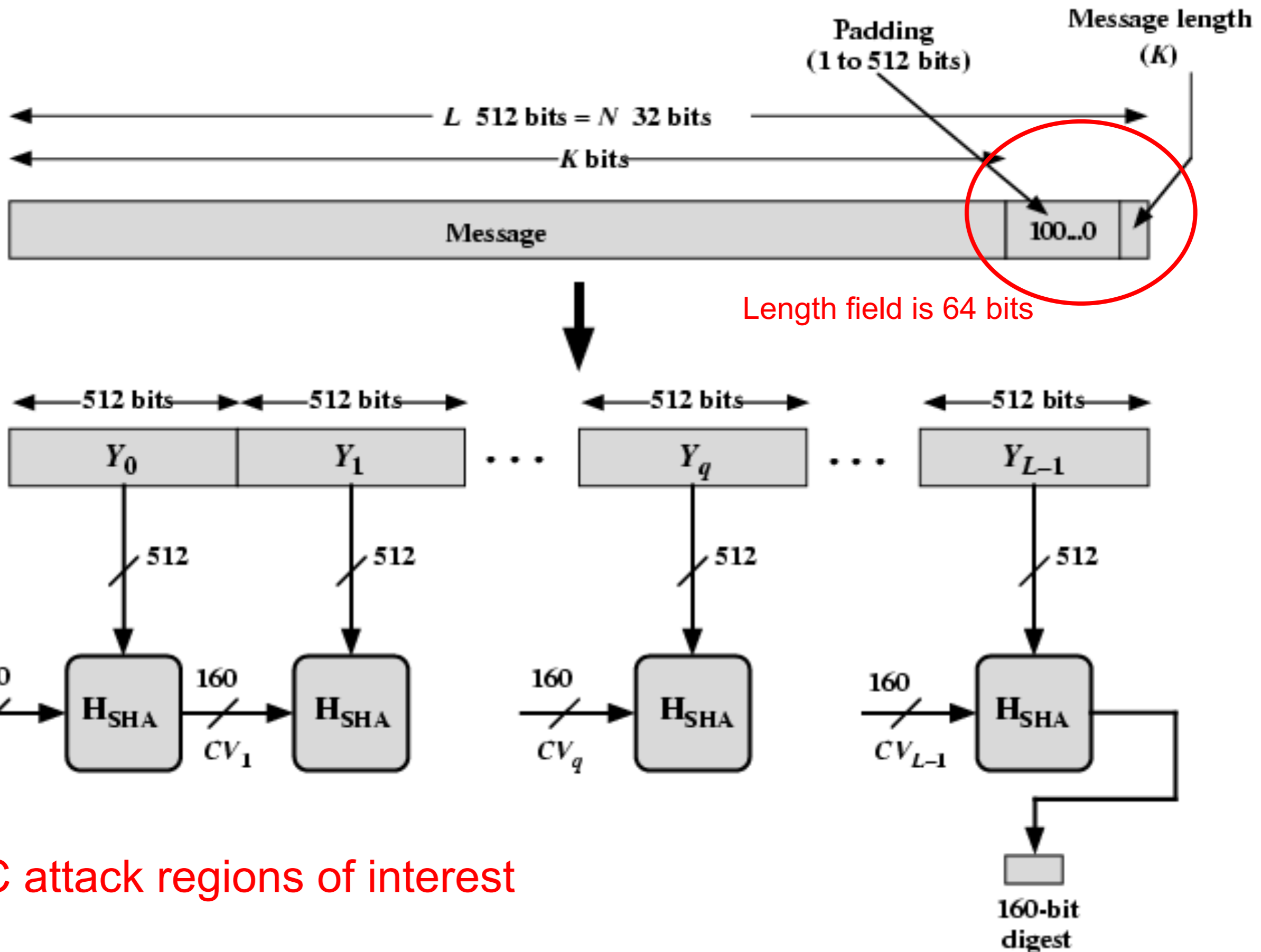
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- Alice and Bob share a key  $K$
- Alice sends message  $M1$  to Bob such that Bob knows it came from Alice
  - Alice computes  $H(K || M1) = mac1$
  - Alice sends  $M1$  and  $mac1$  to Bob
- Bob verifies the message
  - Bob computes  $H(K || M1) = mac2$  and compares it to  $mac1$ . If they match, the message came from Alice.
  - Or did it????

# Message Extension Attack

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- Mallory can intercept a plaintext message and a mac.
- Mallory “extends” the message – adds new material to the end of the message
- She modifies the mac without knowing the key. She needs to know the length of the key.
- She replaces the message and mac with the extended message and new mac and forwards it along
- Bob receives the modified message and mac, and it passes his verification step. He believes it came from Alice!
- See Project 3, resources on that page



**Figure 3.4 Message Digest Generation Using SHA-1**

# Message Extension Attack

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- Extend the message by adding block  $Y_{L+1}$
- If this block is now put through the  $H_{SHA}$  function, along with the output of the previous  $H_{SHA}$  calculation, what do you get?